# Logic and Probabilistic Models of Belief Change 

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- Computer science: updating databases (Doyle 1979 and Fagin et al. 1983)
- Philosophy (epistemology/philosophy of science):
- scientific theory change and revisions of probability assignments;
- belief change (Levi 1977, 1980, Harper 1977) and its rationality.

Carlos Alchourrón, Peter Gärdenfors, and David Makinson.
1985 paper in the Journal of Symbolic Logic.
Starting point of belief revision theory.
C. Alchourrón, P. Gärdenfors and D. Makinson. On the logic of theory change: Partial meet contraction and revision functions. Journal of Symbolic Logic, 50, 510-530, 1985.

## Belief Change

Consider the following beliefs of a rational agent:
$p_{1}$ All Europeans swans are white.
$p_{2}$ The bird caught in the trap is a swan.
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Thus, the agent believes:
$q$ The bird caught in the trap is white.
Now suppose the rational agent-for example, You-learn that the bird caught in the trap is black $(\neg q)$.

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There are several logically consistent ways to incorporate $\neg q$ !

## Belief Change

What extralogical factors serve to determine what beliefs to give up and what beliefs to retain?

## Belief Change

Belief revision is a matter of choice, and the choices are to be made in such a way that:

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2. It is simple; and
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Research has relied on the following related guiding ideas:

1. When accepting a new piece of information, an agent should aim at a minimal change of his old beliefs.
2. If there are different ways to effect a belief change, the agent should give up those beliefs which are least entrenched.

## The Theory of Belief Revision

C. Alchourrón, P. Gärdenfors and D. Makinson. On the logic of theory change: Partial meet contraction and revision functions. Journal of Symbolic Logic, 50, 510-530, 1985.

Hans Rott. Change, Choice and Inference: A Study of Belief Revision and Nonmonotonic Reasoning. Oxford University Press, 2001.
A.P. Pedersen and H. Arló-Costa. "Belief Revision.". In Continuum Companion to Philosophical Logic. Continuum Press, 2011.

## Epistemic States

- Belief sets
- (Ellis's belief systems)
- Possible worlds models
- (Doyle's truth maintenance systems)
- Spohn's generalized possible worlds model
- Bayesian models
- Generalized Bayesian models
- (Johnson-Laird mental models)
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J. Halpern. Reasoning about uncertainty. The MIT Press, 2003.


## AGM

belief := sentence (in some formal language)
beliefs of an agent := a set of such sentences (belief set)

## Language of Beliefs in AGM:

propositional logic: propositions p,q,r,...
connectives: negation ( $\neg$ ), conjunction ( $\wedge$ ), disjunction ( $\vee$ ), implication $(\rightarrow)$, and equivalence $(\leftrightarrow)$.

1. Belief sets should be consistent
2. Belief sets should be closed under logical consequence

## Classical Consequence

For any set $A$ of sentences, $C n(A)$ is the set of logical consequences of $A$. $C n$ is a function from sets of sentences to sets of sentences that satisfies the following three conditions:

- $A \subseteq C n(A)$ (inclusion);
- If $A \subseteq B$, then $C n(A) \subseteq C n(B)$ (monotony);
- $C n(A)=C n(C n(A))$ (idempotence)

If $p$ can be derived from $A$ by classical propositional logic, then $p \in C n(A)$.
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- Belief bases instead of belief sets. $B_{1}=\{p, p \leftrightarrow q\}$, $B_{2}=\{p, q\} . C n\left(B_{1}\right)=C n\left(B_{2}\right)$. What happens when we receive the evidence that $\neg p$ ?


## $K \circ \varphi$



Initial set of beliefs



## Minimal Change

When accepting a new piece of information, an agent should aim at a minimal change of his old beliefs.
"The criterion of informational economy demands that as few beliefs as possible be given up so that the change is in some sense a minimal change of $K$ to accommodate for $A^{\prime \prime}$
(Gardenfors 1988, p. 53).

## Keep the Most Entrenched Beliefs

"A hallmark of the AGM theory is its commitment to the principle of informational economy: beliefs are only given up when there are no less entrenched candidates.... If one of two beliefs must be retracted in order to accommodate some new fact, the less entrenched belief will be relinquished, while the more entrenched persists"
(Boutilier 1996, pp. 264-265).

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2. Contraction: $K \dot{\perp}$; $\varphi$ is removed from $K$ given a new belief set $K^{\prime}$
3. Revision: $K * \varphi ; \varphi$ is added and other things are removed, so that the resulting new belief set $K^{\prime}$ is consistent.

## Contraction Postulates

(C1) $K \doteq \alpha$ is deductively closed
(C2) $K \doteq \alpha \subseteq K$
(C3) If $\alpha \notin K$ or $\vdash \alpha$ then $K \dot{\lrcorner}=K$
(C4) If $\forall \alpha$, then $\alpha \notin K \doteq \alpha$
(C5) If $\vdash \alpha \leftrightarrow \beta$, then $K \dot{\lrcorner}=K \doteq \beta$
$(C 6) \quad K \subseteq C n((K \dot{\bullet}) \cup\{\alpha\})$

## Levi Identity

$$
K * \varphi=(K \dot{-} \varphi)+\varphi
$$

## AGM Postulates

AGM 1: $K * \varphi$ is deductively closed
AGM 2: $\varphi \in K * \varphi$
AGM 3: $K * \varphi \subseteq \operatorname{Cn}(K \cup\{\varphi\})$
AGM 4: If $\neg \varphi \notin K$ then $K * \varphi=C n(K \cup\{\varphi\})$
AGM 5: $K * \varphi$ is inconsistent only if $\varphi$ is inconsistent
AGM 6: If $\varphi$ and $\psi$ are logically equivalent then $K * \varphi=K * \psi$
AGM 7: $K *(\varphi \wedge \psi) \subseteq C n(K * \varphi \cup\{\psi\})$
AGM 8: if $\neg \psi \notin K * \varphi$ then $\operatorname{Cn}(K * \varphi \cup\{\psi\}) \subseteq K *(\varphi \wedge \psi)$

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## Rott's Counterexample

AGM 7: $K *(\varphi \wedge \psi) \subseteq C n(K * \varphi \cup\{\psi\})$

AGM 8: if $\neg \psi \notin K * \varphi$ then $\operatorname{Cn}(K * \varphi \cup\{\psi\}) \subseteq K *(\varphi \wedge \psi)$

So, if $\psi \in \operatorname{Cn}(\{\varphi\})$, then $K * \varphi=\operatorname{Cn}(K * \varphi \cup\{\psi\})$

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There is an appointment to be made in a philosophy department. The position is a metaphysics position, and there are three main candidates: Andrew, Becker and Cortez.

1. Andrew is clearly the best metaphysician, but is weak in logic.
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Scenario 1: Paul is told by the dean, that the chosen candidate is either Andrew or Becker. Since Andrew is clearly the better metaphysician of the two, Paul concludes that the winning candidate will be Andrew.

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Scenario 2: Paul is told by the dean that the chosen candidate is either Andrew, Becker or Cortez.

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Scenario 2: Paul is told by the dean that the chosen candidate is either Andrew, Becker or Cortez.
" This piece of information sets off a rather subtle line of reasoning. Knowing that Cortez is a splendid logician, but that he can hardly be called a metaphysician, Paul comes to realize that his background assumption that expertise in the field advertised is the decisive criterion for the appointment cannot be upheld. Apparently, competence in logic is regarded as a considerable asset by the selection committee." Paul concludes Becker will be hired.
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(Stalnaker, 204)

## Counterexamples to Recovery

$$
K \subseteq C n((K \dot{-}) \cup\{\alpha\})
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While reading a book about Cleopatra I learned that she had both a son and a daughter. I therefore believe both that Cleopatra had a son ( $s$ ) and Cleopatra had a daughter (d).

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(Hansson, 1991)

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(Hansson, 1996)

## Evaluating counterexamples

. . . information about how I learn some of the things I learn, about the sources of my information, or about what I believe about what I believe and don't believe. If the story we tell in an example makes certain information about any of these things relevant, then it needs to be included in a proper model of the story, if it is to play the right role in the evaluation of the abstract principles of the model.

Robert Stalnaker. Iterated Belief Revision. Erkenntnis 70, pp. 189-209, 2009.

Let $K$ be a belief set and $\varphi$ a formula.
$K \perp \varphi$ is the remainder set of $K$.
$A \in K \perp \varphi$ iff

1. $A \subseteq K$
2. $\varphi \notin C n(A)$
3. There is no $B$ such that $A \subset B \subseteq K$ and $\varphi \notin C n(B)$.

- $K \perp \alpha=\{K\}$ iff $\neg \alpha \notin C n(K)$
- $K \perp \alpha=\emptyset$ iff $\alpha \in C n(\emptyset)$
- If $K^{\prime} \subseteq K$ and $\alpha \notin C n\left(K^{\prime}\right)$ then there is some $T$ such that $K^{\prime} \subseteq T \in K \perp \alpha$.

A selection function $\gamma$ for $K$ is a function on $K \perp \alpha$ such that:

- If $K \perp \alpha \neq \emptyset$, then $\gamma(K \perp \alpha) \subseteq K \perp \alpha$ and $\gamma(K \perp \alpha) \neq \emptyset$
- If $K \perp \alpha=\emptyset$, then $\gamma(K \perp \alpha)=\{K\}$

Let $K$ be a set of formulas. A function - is a partial meet contraction for $K$ if there is a selection function $\gamma$ for $K$ such that for all formula $\alpha$ :

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K \dot{-}=\bigcap \gamma(K \perp \alpha)
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Then $K * \alpha=\operatorname{Cn}(\bigcap \gamma(K \perp \neg \alpha) \cup\{\alpha\})$

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Then $K * \alpha=\operatorname{Cn}(\bigcap \gamma(K \perp \neg \alpha) \cup\{\alpha\})$

- $\gamma$ selects exactly one element of $K \perp \alpha$ (maxichoice contraction)
- $\gamma$ selects the entire set $K \perp \alpha$ (full meet contraction)


## AGM Postulates

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Theorem (AGM 1985). Let $K$ be a belief set and let $*$ be a function on $\mathcal{L}$. Then

- The function $*$ is a partial meet revision for $K$ if and only if it satisfies the postulates AGM1 - AGM6
- The function $*$ is a transitively relational partial meet revision for $K$ if and only if it satisfies AGM1 - AGM8.


## Belief Revision: The Semantic View

A. Grove. Two modelings for theory change. Journal of Philosophical Logic, 17, pgs. 157-170, 1988.

EP. Dynamic Epistemic Logic II: Logics of information change. Philosophy Compass, Vol. 8, Iss. 9, pgs. 815-833, 2013.


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- The agent's (hard) information (i.e., the states consistent with what the agent knows)

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- The states consistent with what the agent knows with a distinguished state (the "actual world")
- Each state is associated with a propositional valuation for the underlying propositional language

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## Sphere Models

Let $W$ be a set of states, A set $\mathcal{F} \subseteq \wp(W)$ is called a system of spheres provided:

- For each $S, S^{\prime} \in \mathcal{F}$, either $S \subseteq S^{\prime}$ or $S^{\prime} \subseteq S$
- For any $P \subseteq W$ there is a smallest $S \in \mathcal{F}$ (according to the subset relation) such that $P \cap S \neq \emptyset$
- The spheres are non-empty $\bigcap \mathcal{F} \neq \emptyset$ and cover the entire information cell $\bigcup \mathcal{F}=W$

Let $\mathcal{F}$ be a system of spheres on $W$ : for $w, v \in W$, let

$$
w \preceq_{\mathcal{F}} v \text { iff for all } S \in \mathcal{F} \text {, if } v \in S \text { then } w \in S
$$

Then, $\preceq_{\mathcal{F}}$ is reflexive, transitive, and well-founded.
$w \preceq_{\mathcal{F}} v$ means that no matter what the agent learns in the future, as long as world $v$ is still consistent with his beliefs and $w$ is still epistemically possible, then $w$ is also consistent with his beliefs.

## Belief Revision via Plausibility

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- $w_{1} \prec w_{3}\left(w_{1} \preceq w_{3}\right.$ and $\left.w_{3} \npreceq w_{1}\right)$
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- $w_{2} \prec w_{3}\left(w_{2} \preceq w_{3}\right.$ and $\left.w_{3} \npreceq w_{2}\right)$
- $\left\{w_{1}, w_{2}\right\} \subseteq \operatorname{Min}_{\preceq}\left(\left[w_{i}\right]\right)$



## Belief Revision via Plausibility



Conditional Belief: $B^{\varphi} \psi$

## Belief Revision via Plausibility



Conditional Belief: $B^{\varphi} \psi$

$$
\operatorname{Min}_{\preceq}\left(\llbracket \varphi \rrbracket_{\mathcal{M}}\right) \subseteq \llbracket \psi \rrbracket_{\mathcal{M}}
$$

## Belief Revision via Plausibility



Incorporate the new information $\varphi$

## Belief Revision via Plausibility



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Radical Upgrade: Information from a strongly trusted source $(\Uparrow \varphi): A \prec_{i} B \prec_{i} C \prec_{i} D \prec_{i} E$

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Complete vs. incomplete belief sets:
$K=C n(\{p \vee q\})$ vs. $K=C n(\{p \vee q, p, q\})$
Revising by $\neg p(K * \neg p)$ vs. Updating by $\neg p(K \diamond \neg p)$
H. Katsuno and A. O. Mendelzon. Propositional knowledge base revision and minimal change. Artificial Intelligence, 52, pp. 263-294 (1991).

## KM Postulates

KM 1: $K \diamond \varphi=C n(K \diamond \varphi)$
KM 2: $\varphi \in K \diamond \varphi$
KM 3: If $\varphi \in K$ then $K \diamond \varphi=K$
KM 4: $K \diamond \varphi$ is inconsistent iff $\varphi$ is inconsistent
KM 5: If $\varphi$ and $\psi$ are logically equivalent then $K \diamond \varphi=K \diamond \psi$
KM 6: $K \diamond(\varphi \wedge \psi) \subseteq C n(K \diamond \varphi \cup\{\psi\})$
KM 7: If $\psi \in K \diamond \varphi$ and $\varphi \in K \diamond \psi$ then $K \diamond \varphi=K \diamond \psi$
KM 8: If $K$ is complete then $K \diamond(\varphi \wedge \psi) \subseteq K \diamond \varphi \cap K \diamond \psi$
KM 9: $K \diamond \varphi=\bigcap_{M \in \operatorname{Comp}(K)} M \diamond \varphi$, where $\operatorname{Comp}(K)$ is the class of all complete theories containing $K$.

## Updating and Revising

$$
K \diamond \varphi=\bigcap_{M \in \operatorname{Comp}(K)} M * \varphi
$$

H. Katsuno and A. O. Mendelzon. On the difference between updating a knowledge base and revising it. Belief Revision, P. Gärdenfors (ed.), pp 182-203 (1992).

